#### Conflict resolution: Proper inclusion v. overlap

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## (conflict = competition)

here, competition between generalizations over (phonological) strings

## the point



- Why? What's so special about proper inclusion?
- I argue that the celebrated distinction between proper inclusion and overlap is a spurious one.
  - All that matters is *conflict*, and how it is resolved.

See my 2013 monograph for this same point, embedded in a larger discussion of blocking, complementarity, and the principles that are proposed to regulate these.



## SPE rules and order

In rule-based generative phonology, generalizations are expressed as serially-ordered rewrite rules.



Kenstowicz & Kisseberth (1977, 1979)

Kiparsky (1968)

#### $V \longrightarrow [+stress] / = C_0((\breve{V}C_0^1)VC_0) \#$

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### $V \longrightarrow [+stress] / = C_0 \breve{V} C_0^1 V C_0 \#$

Stress the antepenultimate vowel if there is one and if the penultimate vowel is short and in an open syllable (i.e. the penultimate syllable is light).

#### $V \longrightarrow [+stress] / = C_0((\breve{V}C_0^1)VC_0) \#$

 $V \longrightarrow [+stress] / - C_0 V C_0 \#$ 

Otherwise, stress the penultimate vowel if there is one.

#### $V \longrightarrow [+stress] / = C_0((\breve{V}C_0^1)VC_0) \#$

 $V \longrightarrow [+stress] / - C_0 #$ 

 $\rightarrow$  [+stress] /  $- C_0 \check{V} C_0^1 V C_0 \#$  $\rightarrow$  [+stress] /  $- C_0 V C_0 \#$  $\rightarrow$  [+stress] / — C<sub>0</sub># If application of such rules were pa-trí-cí-á conjunctive rather than disjunctive, there could be as many as three

Chomsky & Halle (1968)

stresses assigned to one word.

 $V \longrightarrow [+stress] / = C_0((\breve{V}C_0^1)VC_0)\#$  $V \longrightarrow [+stress] / = C_0\breve{V}C_0^1VC_0\#$  $V \longrightarrow [+stress] / = C_0VC_0\#$  $V \longrightarrow [+stress] / = C_0\#$ 

Note the proper inclusion relationships among these strings, capitalized upon by the parenthesis notation

## Metrical stress theory

- Final syllable extrametricality (modulo exhaustivity).
- Assign a bimoraic trochee at the right edge.

pa-(trí-ci)-(a) (ré-fi)-〈cit〉 re-(fḗ)-<cit> re-(féc)-<tus> (rế) (méns)

Hayes (1981, 1995)

## Conflict in SPE

Actual conflict between rewrite rules arises under two conditions: mutual feeding and mutual bleeding.



Pullum (1976)

Kiparsky (1971)

### Conflict in SPE



Neither of these types of interactions appears to require anything other than ordering. And yet...

#### Elsewhere Condition

Two rules of the form

 $A \rightarrow B/P \_ Q$  $C \rightarrow D/R \_ S$ 

are disjunctively ordered iff:

- A. the set of strings that fit *PAQ* is a subset of the set of strings that fit *RCS*, and
- B. the structural changes of the two rules are incompatible.

#### Proper Inclusion Precedence Principle

"incompatible structural changes" =  $X \rightarrow Y vs. Y \rightarrow X$ 

the Elsewhere Condition is thus a response to issues involving cases of mutual feeding — it *prevents* Duke of York derivations

A properly includes the structural description of B.

Kiparsky (1973)

Koutsoudas et al. (1974)

#### Eleowhore Condition

"For all the cases of proper inclusion precedence considered here, the related rules are intrinsically disjunctive, since application of either rule yields a representation that fails to satisfy the structural description of the other." (fn. 7, p. 9)

the Proper Inclusion Precedence Principle is thus a response to issues involving cases of mutual bleeding to predict the order of rules in a Duke of Earl relationship

of the two rules are incompatible.

Kiparsky (1973)

#### Proper Inclusion Precedence Principle

For any representation R, which meets the structural description of each of two rules A and B, A takes applicational precedence over B with respect to R iff the structural description of A properly includes the structural description of B.

Koutsoudas et al. (1974)

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#### English lengthening & shortening

- *CiV* Lengthening:  $V \rightarrow \overline{V} / (' \_ C i) V$ 
  - e.g. re('mēdi)(al), ('rādi)(al), me('lōdi)(ous)...
- Trisyllabic Shortening:  $V \rightarrow \breve{V} / (' \_ C_0 V)$ 
  - e.g. ('rĕme) $\langle dy \rangle$ , ('rădi) $\langle cal \rangle$ , ('mĕlo) $\langle dy \rangle$ ...



Kenstowicz (1994)

#### English lengthening & shortening

| ✤ = blocking by EC  | ('rådi)(al) | ('rådi)(cal) |
|---|-------------|--------------|
| $Lengthening V \longrightarrow \overline{V} / (' \_ C i) V$ | ('rādi)〈al〉 |              |
| Shortening $V \rightarrow \breve{V} / (' \_ C_0 V)$         | *           | ('rădi)(cal) |

Kenstowicz (1994)

#### English lengthening & shortening

| Just to avoid this?   | ('rådi)(al) | ('rådi)(cal) |
|---|-------------|--------------|
| Shortening $V \rightarrow \breve{V} / (' \_ C_0 V)$         | ('rădi)(al) | ('rădi)(cal) |
| $Lengthening V \longrightarrow \overline{V} / (' \_ C i) V$ | ('rādi)(al) |              |

#### Disjunctive application is "maximized".

Chomsky (1967: 124-125), Chomsky & Halle (1968: 63)

"[C]ertain natural economy conditions" require that there be "no 'superfluous steps' in derivations".

Chomsky (1995: 220), Halle & Idsardi (1998: 1)

## Nootka / Nuuchahnulth

#### labialization & delabialization

| <i>Overlap requires<br/>Duke of York!</i>                      | muq              | ħaju-qi               | łaːkʷ-∫itł |
|--|------------------|-----------------------|------------|
| $Labialization \\ [dors] \rightarrow [+rd] / [+rd]$            | muq <sup>w</sup> | ħaju-q <sup>∞</sup> i |            |
| $\frac{Delabialization}{[dors] \rightarrow [-rd] / \{\sigma}}$ | muq              |                       | łaːk-∫itł  |

## Whence proper inclusion?

- Proper inclusion is the one subcase of overlap for which there is only one truly possible order.
  - General > Specific allows Specific to apply,
  - Specific > General occults Specific.
- Proper inclusion is asymmetrically complete; unique among forms of overlap in that it can be non-arbitrarily used to determine which of two conflicting rules is blocked.

#### English' lengthening & shortening

| Rules reversed  | ('rådi)(al)       | ('rådi)(cal) |
|---|-------------------|--------------|
| $Lengthening V \longrightarrow \overline{V} / (' \_ C i) V$ | ('rādi)〈al〉       |              |
| Shortening $V \rightarrow \breve{V} / (' \_ C_0 V)$         | ('rădi) <al></al> | ('rădi)(cal) |

## Nootka / Nuuchahnulth'

labialization & delabialization

| Rules reversed  | muq | ħaju-qi  | łaːkʷ-∫itł |
|---|-----|----------|------------|
| $Delabialization \\ [dors] \longrightarrow [-rd] / \{\sigma}$ | muq |          | łaːk-∫itł  |
| $Labialization \\ [dors] \longrightarrow [+rd] / [+rd]$       | muq | ħaju-qʷi |            |

# So what counts as conflict?

## English

| $/(\text{'rådi})\langle al \rangle/$     | CiV-Long | Short | F   |
|--|----------|-------|-----|
| a. IS $(r\bar{a}di)\langle al \rangle$   |          | *     | (*) |
| b. $('r \breve{a} di)\langle al \rangle$ | *!       |       | (*) |

| $/(\text{'rådi})\langle \text{cal} \rangle /$ | CiV-Long | Short | F   |
|---|----------|-------|-----|
| a. $(r\bar{a}di)\langle al \rangle$           |          | * !   | (*) |
| b. 🖙 ('rădi)(al)                              |          |       | (*) |

## English'

| $/(\text{'rådi})\langle al \rangle/$ | Short | CiV-Long | F   |
|--------------------------------------|-------|----------|-----|
| a. $(r\bar{a}di)\langle al \rangle$  | *!    |          | (*) |
| b. ☞ ('rădi)⟨al⟩                     |       | *        | (*) |

| $/(\text{'rådi})\langle \text{cal} \rangle /$ | Short | CiV-Long | F   |
|---|-------|----------|-----|
| a. $(r\bar{a}di)\langle al \rangle$           | * !   |          | (*) |
| b. 🖙 ('rădi)(al)                              |       |          | (*) |

### Nootka / Nuuchahnulth

| /muq/               | Delab | Lab | F   |
|---------------------|-------|-----|-----|
| a. 🖙 muq            |       | *   | (*) |
| b. muq <sup>w</sup> | * !   |     | (*) |

| /ħaju-qi/                  | Delab | Lab | F   |
|----------------------------|-------|-----|-----|
| a. ħaju-qi                 |       |     | (*) |
| b. ☞ ħaju-q <sup>w</sup> i |       | * ! | (*) |

### Nootka / Nuuchahnulth'

| /muq/                 | LAB | Delab | F   |
|-----------------------|-----|-------|-----|
| a. muq                | *!  |       | (*) |
| b. ☞ muq <sup>w</sup> |     | *     | (*) |

| /ħaju-qi/                  | LAB | Delab | F   |
|----------------------------|-----|-------|-----|
| a. ħaju-qi                 |     |       | (*) |
| b. ☞ ħaju-q <sup>w</sup> i | *!  |       | (*) |

#### mutual feeding

## 'obliterative bleeding'

Kiparsky (1973)



Kiparsky (1973)

#### Diola Fogny assimilation & deletion

| ✤ = blocking by EC   | ni-gam-gam | na-laŋ-laŋ | let-ku-jaw |
|--|------------|------------|------------|
| Assimilation<br>N $\rightarrow [\alpha pl] / \_ [\alpha pl, -ct]$                            | ni-gaŋ-gam |            |            |
| $\begin{array}{c} \text{Deletion} \\ C \longrightarrow \varnothing \ / \ \_ \ C \end{array}$ | *          | na-la-laŋ  | le-ku-jaw  |

Kiparsky (1973)

# Diola Fogny' assimilation & deletion

| this order   | ni-gam-gam | na-laŋ-laŋ | let-ku-jaw      |
|--|------------|------------|-----------------|
| Assimilation<br>$N \rightarrow [\alpha pl] / \_ [\alpha pl, -ct]$                            | ni-gaŋ-gam |            |                 |
| $\begin{array}{c} \text{Deletion} \\ C \longrightarrow \varnothing \ / \ \_ \ C \end{array}$ | ni-ga-gam  | na-la-laŋ  | le-ku-jaw       |
| 'obliterat<br>bleeding   | ive<br>g'  |            | Kiparsky (1973) |

#### Diola Fogny" assimilation & deletion

| that order   | ni-gam-gam | na-laŋ-laŋ | let-ku-jaw |  |  |
|--|------------|------------|------------|--|--|
| $\begin{array}{c} \text{Deletion} \\ C \longrightarrow \varnothing \ / \ \_ \ C \end{array}$ | ni-ga-gam  | na-la-laŋ  | le-ku-jaw  |  |  |
| Assimilation<br>$N \rightarrow [\alpha pl] / [\alpha pl, -ct]$                               |            |            |            |  |  |
| ʻobliterative<br>bleeding'<br>Kiparsky (1973   |            |            |            |  |  |

## Diola Fogny'(')

| $\left\{ V \begin{array}{c} N \\ \left[ (-)\alpha pl \right] \\ \sigma \end{array} \right]_{\sigma} \left[ \begin{array}{c} -ct \\ \alpha pl \\ \end{array} \right] \right\}$          | Agr(pl)-NC | NoCoda-C | ID(pl) | Max-C |
|--|------------|----------|--------|-------|
| a. $\mathbb{R} \left\{ V \right\}_{\sigma} \begin{bmatrix} -ct \\ \alpha pl \end{bmatrix} \right\}$  |            |          |        | *     |
| b. $\odot \{ V \ N \ [\alpha pl] \ ]_{\sigma} \begin{bmatrix} -ct \\ \alpha pl \end{bmatrix} \}$   |            | * !      | (*)    |       |
| $\left[ \begin{array}{cc} c. & \left\{ V \begin{array}{c} N \\ \left[ -\alpha pl \right] \end{array} \right]_{\sigma} \begin{bmatrix} -ct \\ \alpha pl \end{bmatrix} \right\} \right]$ | *!         | *        | (*)    |       |

#### Elsewhere Condition

Two rules of the form  $A \rightarrow B/P \_ Q$  $C \rightarrow D/R \_ S$ 

are disjunctively ordered iff:

- the set of strings that fit *PAQ* is a subset of the set of strings that fit *RCS*, and
- the structural changes of the two rules are incompatible.

#### Elsewhere Condition

Rules A, B apply disjunctively to a form  $\Phi$  iff

- The structural description of A properly includes that of B.
- The result of applying A to Φ is distinct from the result of applying B to Φ.
  In that case, A is applied first, and if it takes effect, then B is not applied.

Kiparsky (1973)

Kiparsky (1982)

#### Elsewhere Condition

Two rules of the form  $A \rightarrow B / P = Q$  $C \rightarrow D / R = S$ 

are disjunctively ordered iff:

- the set of strings that fit *PAQ* **is a subset of** the set of strings that fit *RCS*, and
- the structural changes of the two rules are **incompatible**.

#### Elsewhere Condition

Rules A, B apply disjunctively to a form Φ iff

- The structural description of A **properly includes** that of B.
- The result of applying A to Φ is distinct from the result of applying B to Φ.

In that case, A is applied first, and if it takes effect, then B is not applied.

Kiparsky (1973)

Kiparsky (1982)

- The result of applying Assimilation is certainly "distinct" from the result of applying Deletion.
- But the result of applying Palatalization is also "distinct" from the result of applying Voicing, and yet we expect them both to apply in this case.



## An alternative for Diola

- Prosodic licensing
  - consonants linked to the onset are licensed,
  - consonants not linked to the onset are deleted.
- Effectively: Deletion only targets unassimilated Cs
  - Assimilation simply bleeds Deletion.

## Diola Fogny

| $\left\{ V \begin{array}{c} N \\ \left[ (-)\alpha \mathrm{pl} \right] \sigma \end{array} \right]_{\sigma} \left[ \begin{array}{c} -\mathrm{ct} \\ \alpha \mathrm{pl} \end{array} \right] \right\}$ | Agr(pl)-NC | NoCoda-C | ID(pl) | Max-C |
|--|------------|----------|--------|-------|
| a. {V ] <sub><math>\sigma</math></sub> [ $-ct$ <sub><math>\alpha</math>pl</sub> }  |            |          |        | * !   |
| b. $\mathbb{R}\left\{ V \mid N \\ \left[ \alpha \mathrm{pl} \right] \right\} = \left\{ \left[ \begin{array}{c} -\mathrm{ct} \\ \alpha \mathrm{pl} \end{array} \right] \right\}$                    |            |          | (*)    |       |
| c. {V N $\begin{bmatrix} -ct \\ \alpha pl \end{bmatrix}_{\sigma} \begin{bmatrix} -ct \\ \alpha pl \end{bmatrix}$ }   | *!         | *        | (*)    |       |

#### Another alternative

| $\left[ \left\{ V \begin{array}{c} N \\ \left[ (-)\alpha pl \right] \right]_{\sigma} \left[ \begin{matrix} -ct \\ \alpha pl \end{matrix} \right] \right\} \right]$                     | Agr(pl)-NC | NoCoda-C | ID(pl)     | Max-C |
|--|------------|----------|------------|-------|
| a. {V ] <sub><math>\sigma</math></sub> [ $-ct$ <sub><math>\alpha</math>pl</sub> }  | * !        |          |            | *     |
| b. $\mathbb{R}\left\{ V \mid N \\ [\alpha pl] \right\}_{\sigma} \begin{bmatrix} -ct \\ \alpha pl \end{bmatrix} \right\}$   |            |          | ;<br>; (*) |       |
| $\left[ \begin{array}{cc} c. & \left\{ V \begin{array}{c} N \\ \left[ -\alpha pl \right] \end{array} \right]_{\sigma} \begin{bmatrix} -ct \\ \alpha pl \end{bmatrix} \right\} \right]$ | * !        | *        | (*)        |       |

Baković (2009)

## recall the point

- I have argued that the celebrated distinction between proper inclusion and overlap is a spurious one.
  - All that matters is *conflict*, and how it is resolved.

Thank you.